Problèmes d'identification dans les graphes

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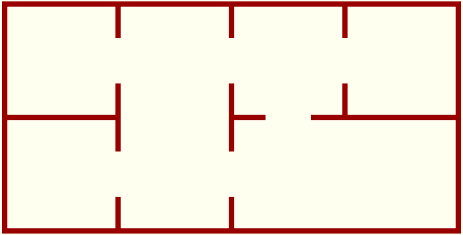
5 juillet 2012

maths à modeler

INSTITUT FOURIER

UNIVERSITÉ DE GRENOBLE

Fire detection in a museum?



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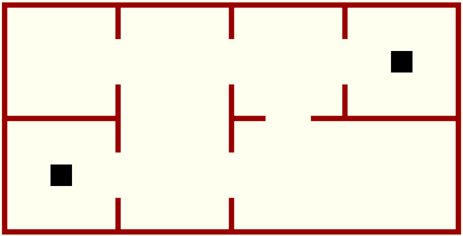

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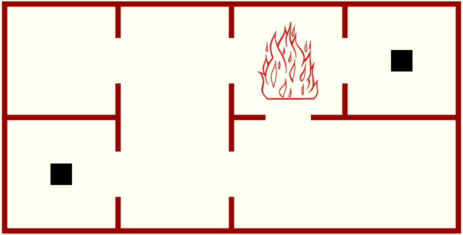

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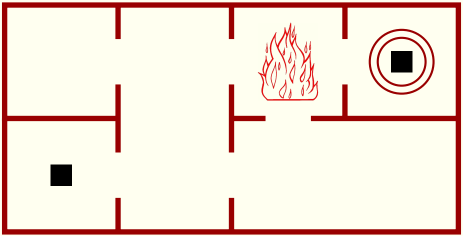

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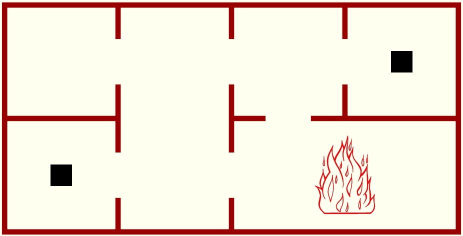

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Fire detection in a museum?

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- Each room must contain a detector or have a detector in a neighboring room.

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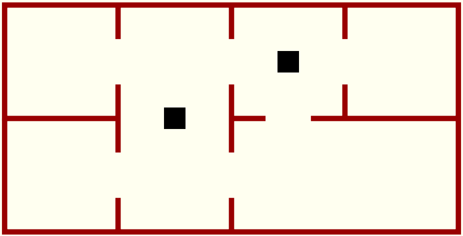

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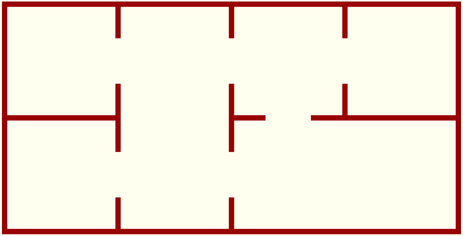
Modelization with a graph



- Vertices V: rooms
- Edges E: between two neighboring rooms

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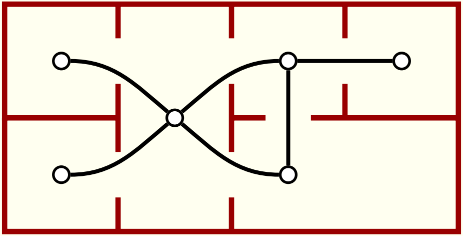

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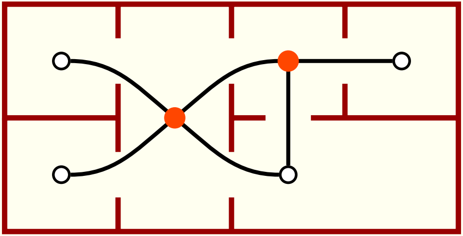

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- Set of detectors = dominating set S:

$$\forall u \in V, N[u] \cap S \neq \emptyset$$

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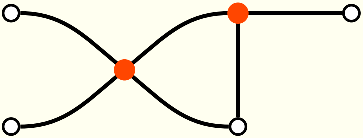

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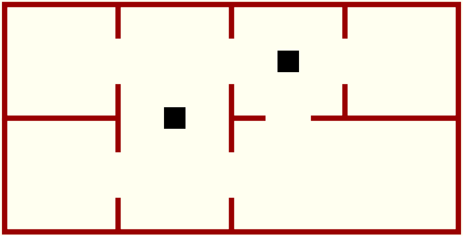

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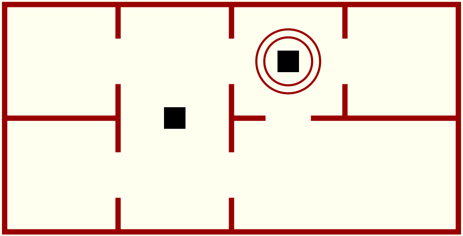

Where is the fire ?

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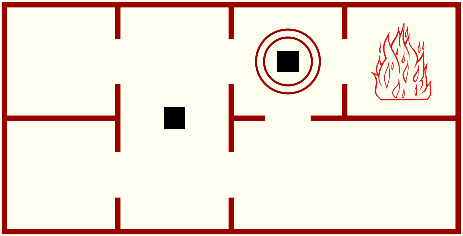

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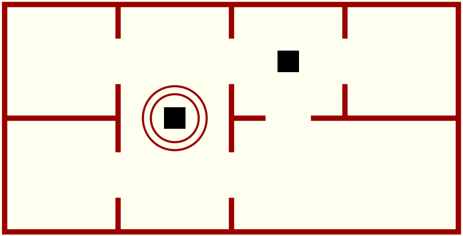

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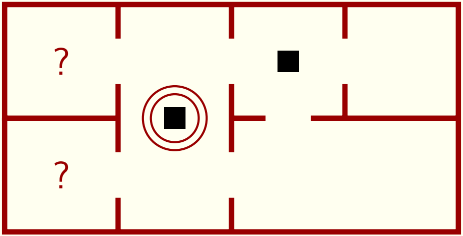
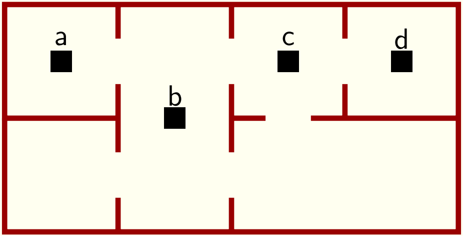


Where is the fire ?

To locate the fire, we need more detectors.

Identifying where is the fire



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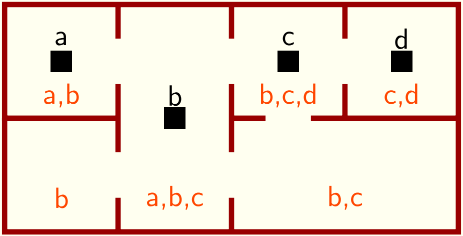
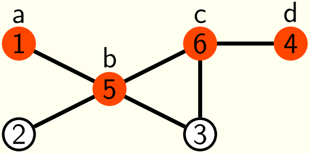


In each room, the set of detectors in the neighborhood is unique.

Modelization with a graph

Identifying code C = subset of vertices which is

- dominating : $\forall u \in V, N[u] \cap C \neq \emptyset$,
- separating : $\forall u, v \in V, N[u] \cap C \neq N[v] \cap C$.



V \ C	a	b	c	d
1	•	•	-	-
2	-	•	-	-
3	-	•	•	-
4	-	-	•	•
5	•	•	•	-
6	-	•	•	•

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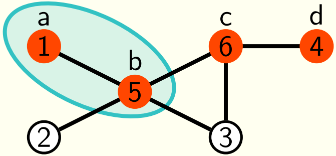

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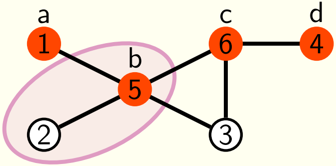

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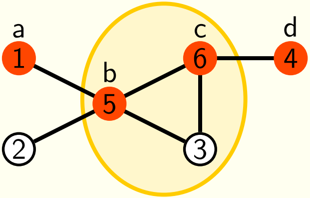

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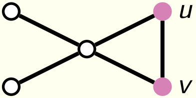
Facts about identifying codes

- Introduced in 1998 by Karpvosky, Chakrabarty and Levitin
- Existence $\Leftrightarrow$ no twins in the graph:



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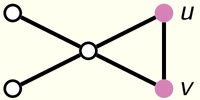

Twins: $N[u] = N[v]$

Given a twin-free graph G, what is the size $\gamma^{ID}(G)$ of minimum identifying code ?

A difficult question...

IDENTIFYING CODE : Given a twin-free graph G and an integer k, is there an identifying code of size k in G?

Proposition Charon, Hudry, Lobstein, 2001

IDENTIFYING CODE is NP-complete.

Outline

1. Bounds and extremal graphs
2. Study in restricted classes of graphs
3. Variation of the definition

Part I

Bounds and extremal graphs

Bounds

$$\log(|V|+1) \leq \gamma^{ID}(G) \leq |V|-1$$

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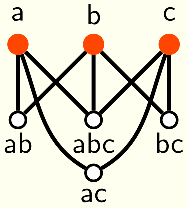
- Tight example:



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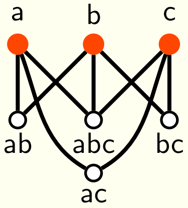

- Complete characterization by Moncel in 2006.

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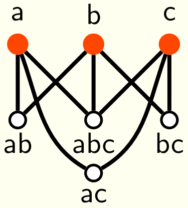

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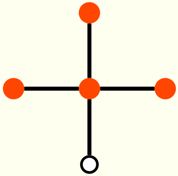
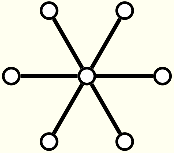
- Bertrand and Gravier, Moncel in 2001.
- Tight example:



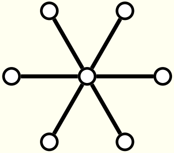
- Complete characterization?

Some tight examples and a conjecture



Stars

Some tight examples and a conjecture



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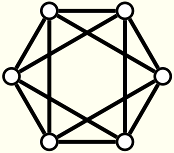
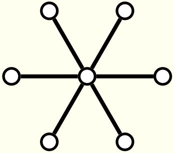


Complete graphs minus maximal matching

Some tight examples and a conjecture



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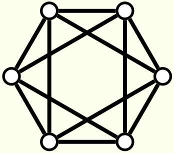


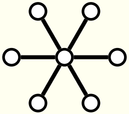
Complete graphs minus maximal matching

Conjecture Charbit, Charon, Cohen, Hudry, Lobstein, 2008

These are the only graphs with $\gamma^{ID} = |V| - 1$.

Characterization of graphs with $\gamma^{ID}(G) = |V| - 1$

(1) Star $K_{1,n}$,



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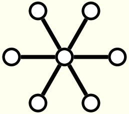

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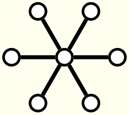

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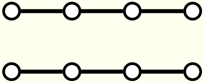

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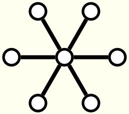

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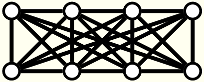

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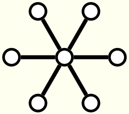

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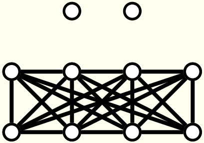

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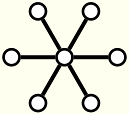

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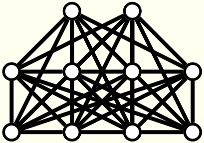

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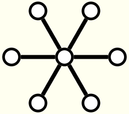

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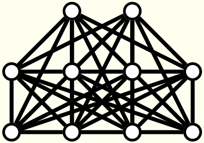

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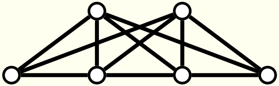

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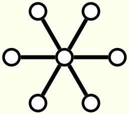

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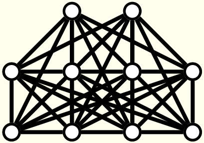

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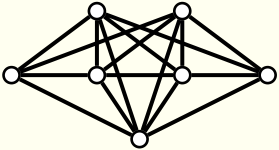

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Theorem Foucaud, Guerrini, Kovše, Naserasr, P., Valicov, 2011

Let G be a connected twin-free graph.

$$\gamma^{ID}(G) = |V| - 1 \Leftrightarrow G \text{ in } (1), (2), (3) \text{ or } (4)$$

Ideas of the proof

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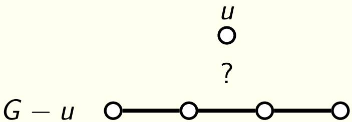
$\Leftarrow$ By induction

$\Rightarrow$ Let G be a minimal counter-example.

- There is $u \in V$ s.t. $G - u$ extremal.
- By minimality, $G - u$ is in (1), (2), (3) or (4).
- We can construct an identifying code of size $|V| - 2$ of G, contradiction.

Ideas of the proof - one small case

Example: $G - u = P_4$



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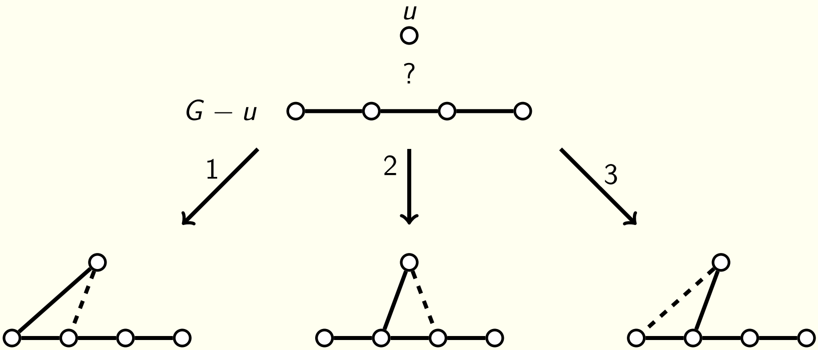

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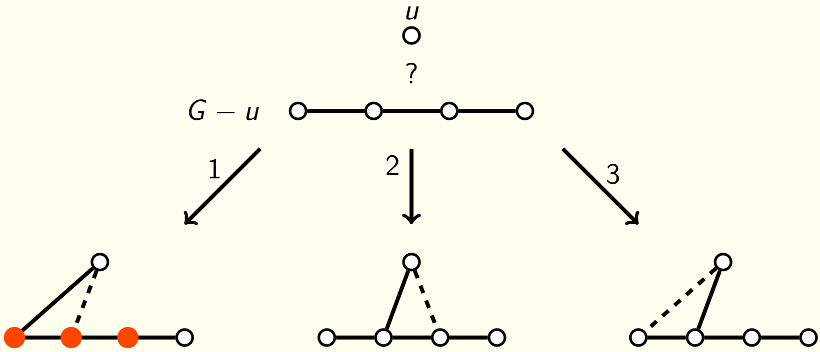

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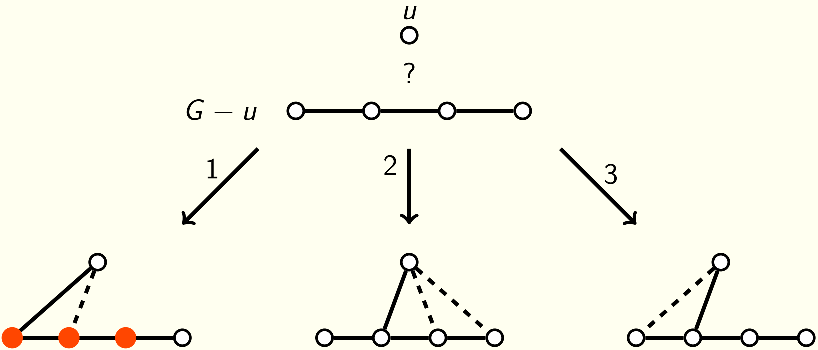

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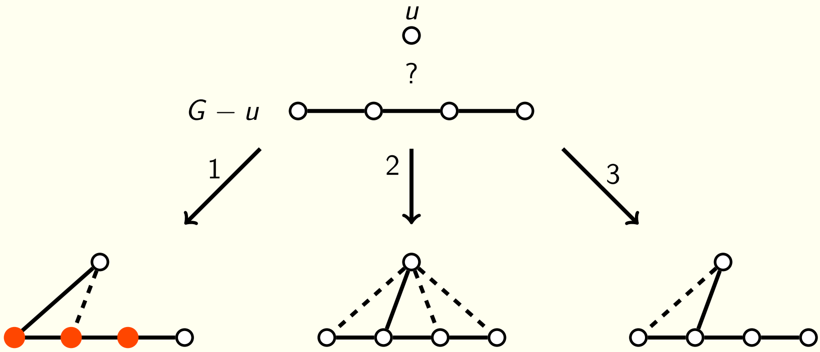

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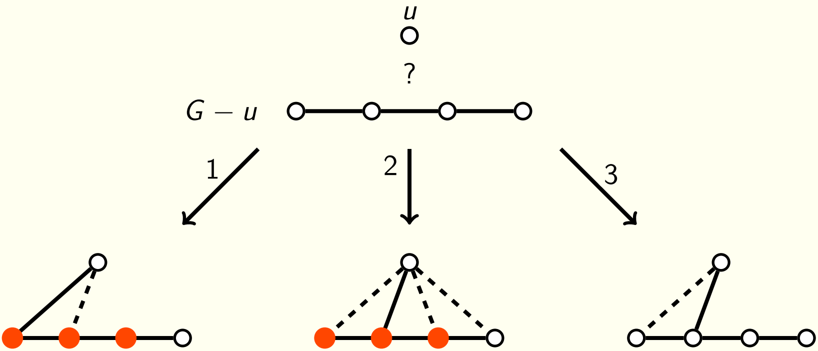

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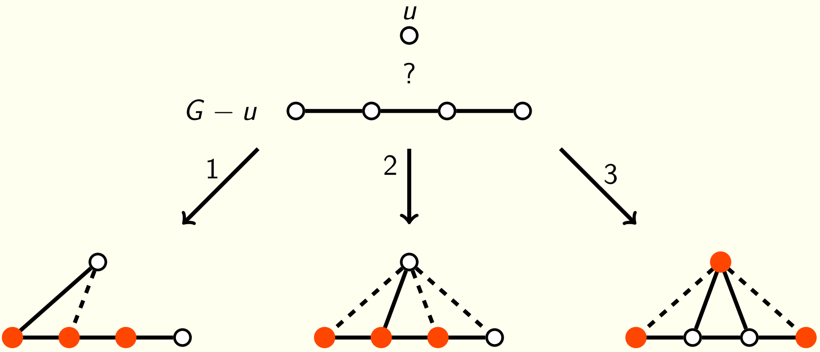

In each case, there is an identifying code of size 3.

Consequence

Corollary

If $\gamma^{ID}(G) = |V| - 1$, G has maximum degree $\Delta \geq |V| - 2$.

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Corollary

If $\gamma^{ID}(G) = |V| - 1$, G has maximum degree $\Delta \geq |V| - 2$.

Upper bound with the maximum degree Δ?

Conjecture Foucaud, Klasing, Kosowski, Raspaud, 2012

$$\gamma^{ID}(G) \leq |V| - \frac{|V|}{\Delta} + O(1).$$

Similar results

Characterization of graphs for which the only IC is the whole vertex set for:

- Infinite non oriented graphs
- Finite digraphs (Foucaud, Naserasr, P., 2012)
- Infinite oriented graphs (Foucaud, Naserasr, P., 2012)

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- Infinite non oriented graphs
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  $\rightarrow$ Consequence for finite graphs: $\gamma^{ID}(G) \leq |V| - \frac{|V|}{\Theta(\Delta^5)}$.
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- Finite digraphs (Foucaud, Naserasr, P., 2012)
  $\rightarrow$ No oriented cycle.

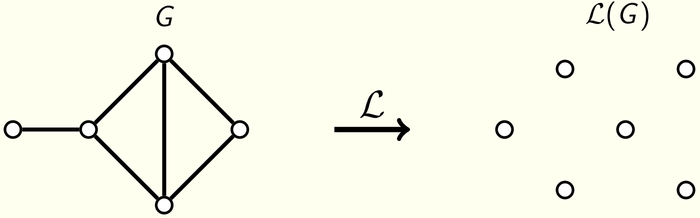
- Infinite oriented graphs (Foucaud, Naserasr, P., 2012)

Part II

Study in a restricted class of graphs:

Line graphs

Identifying code in line graphs



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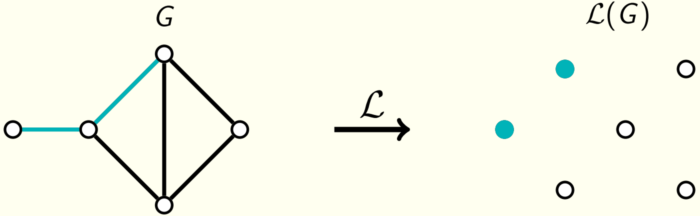

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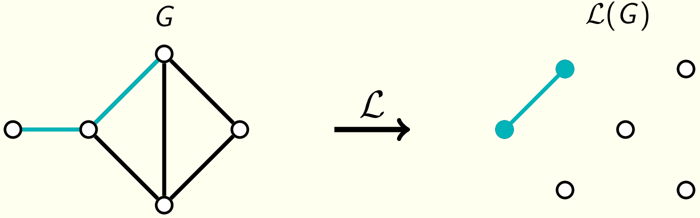

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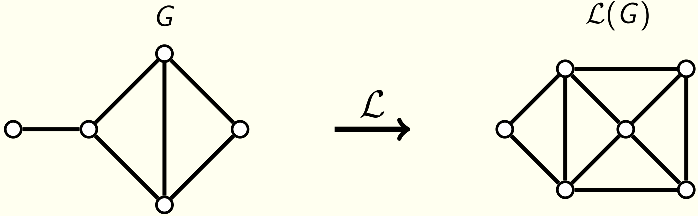

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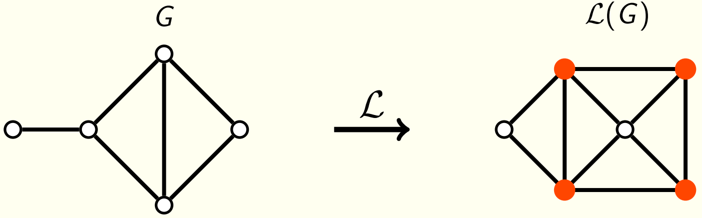
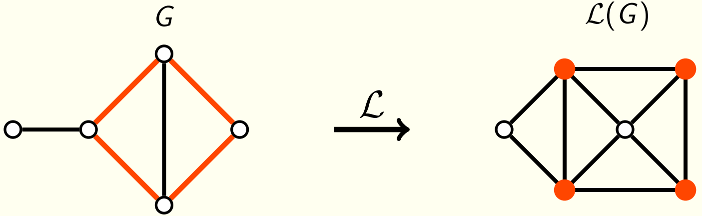


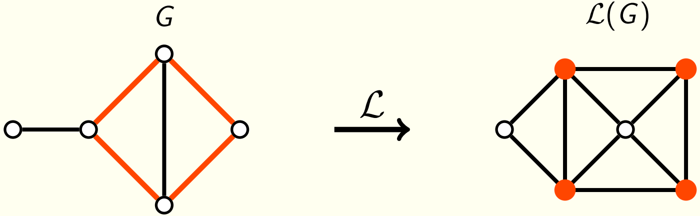
Identifying code

Identifying code in line graphs



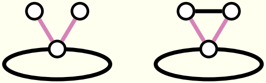
Edge identifying code ⟷ Identifying code

Identifying code in line graphs

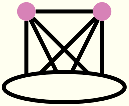


Edge identifying code ⟷ Identifying code

$\gamma^{EID}(G) = \gamma^{ID}(\mathcal{L}(G))$



⟷



Pendant edges ⟷ Twins

Still difficult

EDGE-IDCODE : Given G pendant-free and k, $\gamma^{EID}(G) \leq k$?

Theorem Foucaud, Gravier, Naserasr, P., Valicov, 2012

EDGE-IDCODE is NP-complete even for planar subcubic bipartite graphs with large girth.

Reduction from PLANAR $(\leq 3, 3)$-SAT.

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Reduction from PLANAR $(\leq 3, 3)$-SAT.

Corollary

IDENTIFYING CODE is NP-complete even for perfect planar 3-colorable line graphs with maximum degree 4.

Bounds using the number of vertices

Proposition Foucaud, Gravier, Naserasr, P., Valicov, 2012

$$\frac{1}{2}|V(G)| \leq \gamma^{EID}(G) \leq 2|V(G)| - 3$$

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- Lower Bound: a code must cover $\simeq$ half of vertices.
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- Upper Bound: a minimal code is 2-degenerate.
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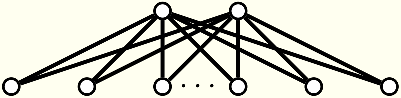
  $\rightarrow$ Infinite family with $\gamma^{EID}(G) = 2|V(G)| - 6$:



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Corollary

EDGE-IDCODE has a polynomial 4-approximation.

- Best polynomial approximation for identifying codes in $\log(|V|)$.
  (Laifenbeld, Trachtenberg, Berger-Wolf, 2006 and Gravier, Klasing, Moncel, 2008)

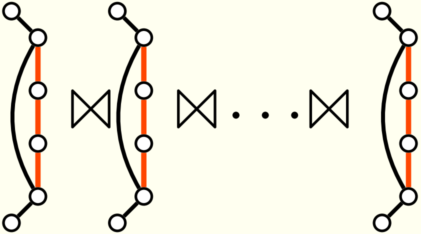
Bounds using the number of edges

Proposition Foucaud, Gravier, Naserasr, P., Valicov, 2012

$$\frac{3}{2\sqrt{2}}\sqrt{|E(G)|} \leq \gamma^{EID}(G) \leq |E(G)| - 1$$

- Upper Bound: from identifying code
- Lower Bound: using the lower bound for vertices

$\rightarrow$ Tight for:



Bounds using the number of edges

Proposition Foucaud, Gravier, Naserasr, P., Valicov, 2012

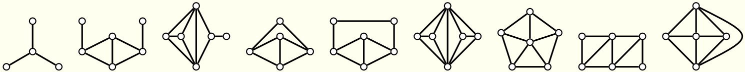
$$\frac{3}{2\sqrt{2}}\sqrt{|E(G)|} \leq \gamma^{EID}(G) \leq |E(G)| - 1$$

Corollary

If G is a line graph, $\gamma^{ID}(G) \geq \Theta(\sqrt{|V|})$

Conclusion for line graphs

- Class of graph for which $\gamma^{ID}(G) \geq \Theta(\sqrt{|V|})$.
- Defined by forbidden induced subgraphs:



- Is the lower bound still true with less restrictions? For other classes defined by forbidden induced subgraphs?
  - $\rightarrow$ False for claw-free graphs.
  - $\rightarrow$ True for interval graphs.

Part III

A variation of identifying code:

Identifying colorings of graphs

Some variations

- Locating-dominating codes
- Resolving sets
- $(r, \leq \ell)$-identifying codes
- Weak and light codes
- Tolerant identifying codes
- Watching systems
- Discriminating codes
- Adaptative identifying codes
- Locating colorings
- ...

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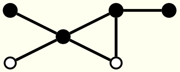
One more:

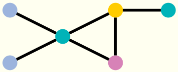
Identifying coloring

Identification with colors

Identifying codes

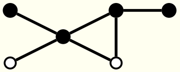


(Proper) graph colorings



Identification with colors

Identifying codes



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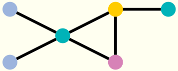
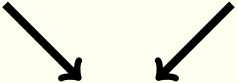




Identifying colorings

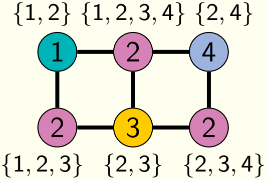
Similar colorings

- Vertex-distinguishing edge-colorings (Harary & Plantholt, 1985)
- Adjacent vertex-distinguishing edge-colorings (Zhang, Liu, Wang, 2002)
- Vertex-distinguishing total colorings (Zhang, Chen, Li, Yao, Lu, Wang, 2005)

Identifying coloring

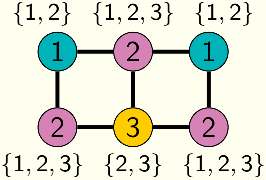
- Vertex coloring $c : V \to \mathbb{N}$
- Vertex identified by the colors in the neighborhood: $c(N[x])$

Global



$c(N[x]) \neq c(N[y])$ for all x,y

Local

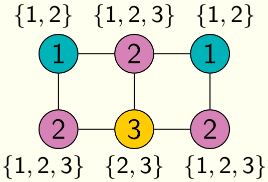


$c(N[x]) \neq c(N[y])$ for $xy \in E$

Definition of locally identifying coloring

A locally identifying coloring (lid-coloring) of G is a coloring c s.t., for each edge xy:

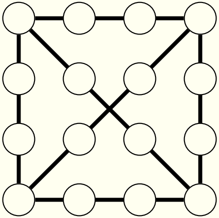
- $c(x) \neq c(y)$ (proper coloring)
- if $N[x] \neq N[y]$, $c(N[x]) \neq c(N[y])$



- $\chi_{lid}(G)$: min. number of colors in a lid-coloring of G.

Link with chromatic number

- A lid-coloring is a proper coloring: $\chi_{lid} \geq \chi$.
- No upper bound with χ.
  $\rightarrow$ complete graph K_k subdivided twice: $\chi_{lid} = k$, $\chi = 3$



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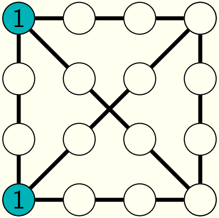

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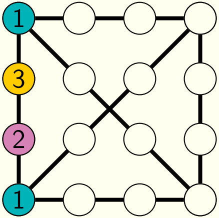

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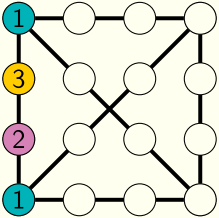
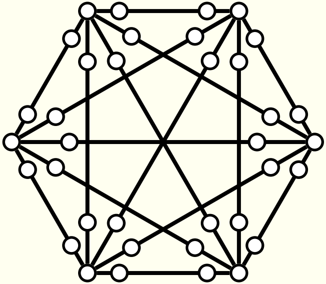


- Not monotone: $\chi_{lid}(P_5) \leq \chi_{lid}(P_4)$

χ_{lid} is not monotone at all



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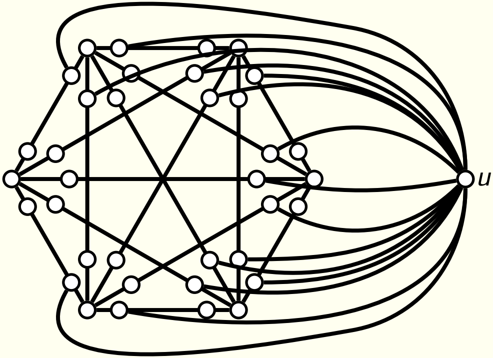
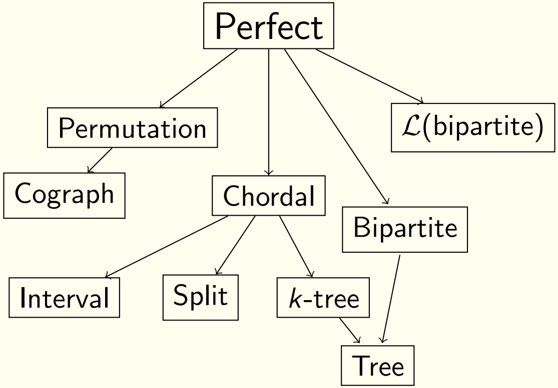
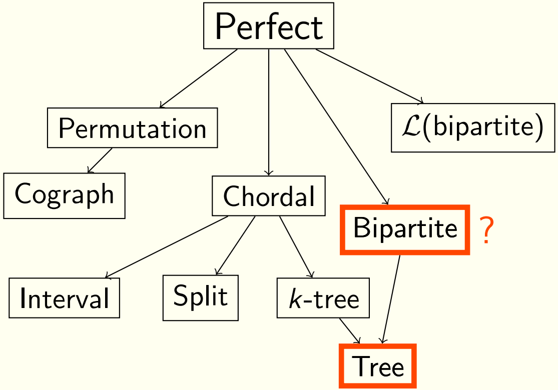


$$\chi_{lid}(G) = 5 \ll k = \chi_{lid}(G - u)$$

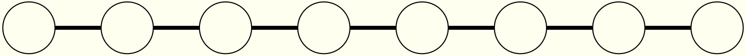
Study in perfect graphs



Study in perfect graphs



Bipartite graphs: the path



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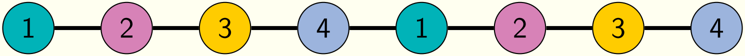

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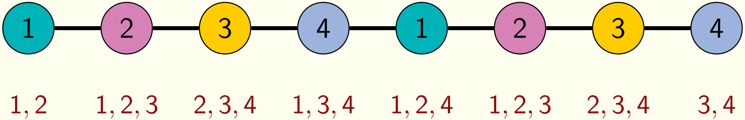
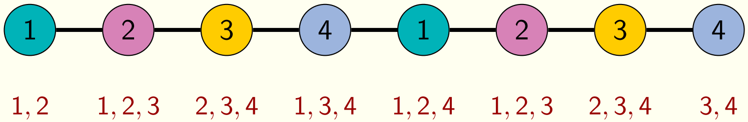
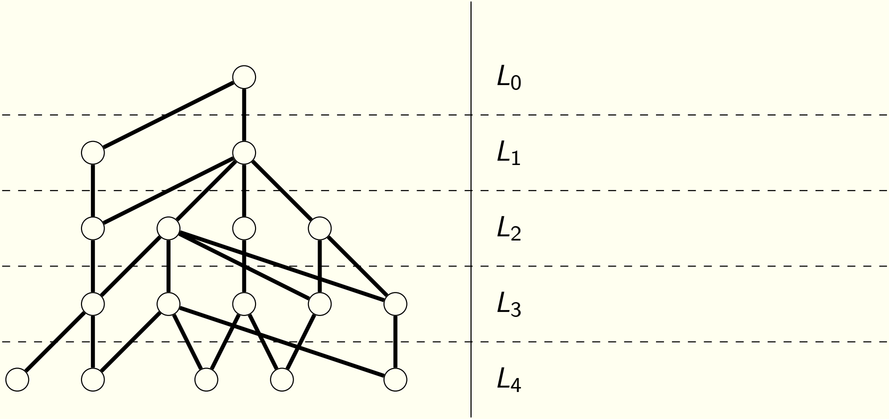


Bipartite graphs: the path



$$\chi_{lid}(P_k) \leq 4$$

Bipartite graphs are 4-lid-colorable



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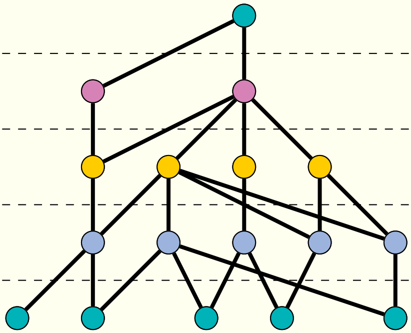
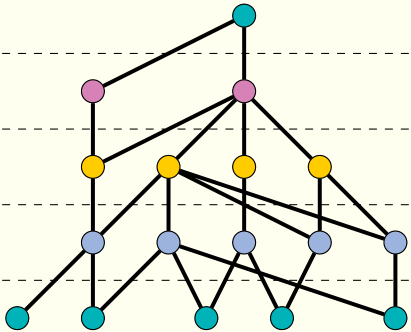


L_0	$\rightarrow$	1	$\{1,2\}$
L_1	$\rightarrow$	2	$\{1,2,3\}$
L_2	$\rightarrow$	3	$\{2,3,4\}$ or $\{2,3\}$
L_3	$\rightarrow$	4	$\{1,3,4\}$ or $\{3,4\}$
L_4	$\rightarrow$	1	$\{1,4\}$

Bipartite graphs are 4-lid-colorable



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If G is bipartite, $\chi_{lid}(G) \leq 4$.

Bipartite graphs

General bounds: $3 \leq \chi_{lid}(B) \leq 4$.

Bipartite graphs

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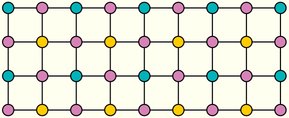
$\chi_{lid}(B) = 3$:



Bipartite graphs

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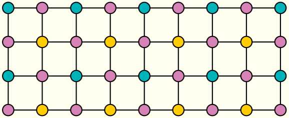

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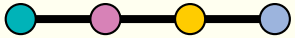

Bipartite graphs

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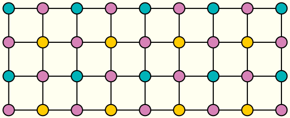

$\leftarrow$? $\rightarrow$

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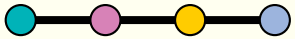
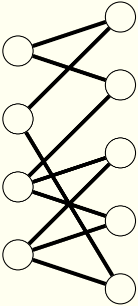


In general... 3-LID-COLORING is NP-complete in bipartite graphs

Link with 2-coloring of hypergraph

Try to color a bipartite graph with 3 colors.



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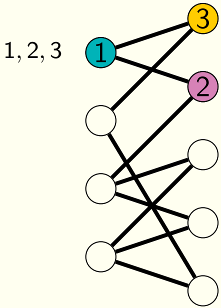

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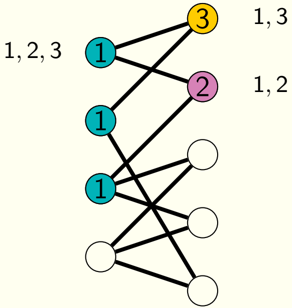

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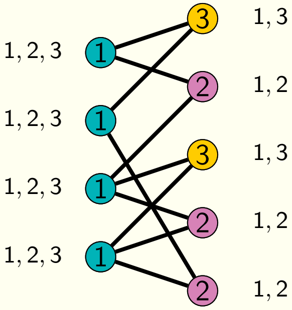

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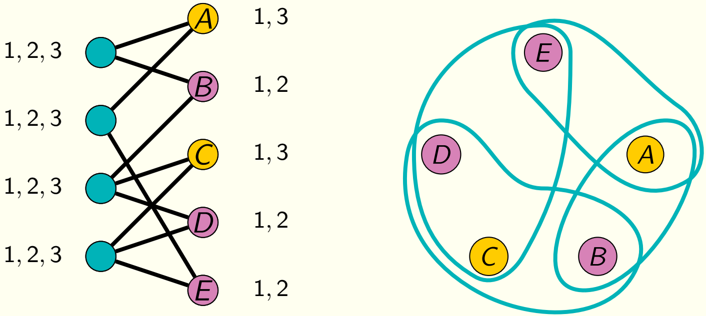
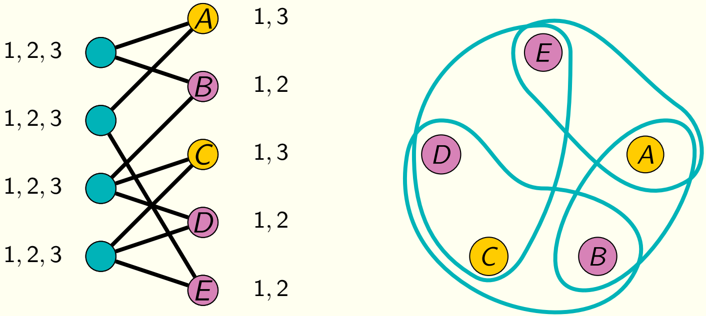


3-lid-coloring in bipartite graph ⇔ 2-coloring in hypergraph

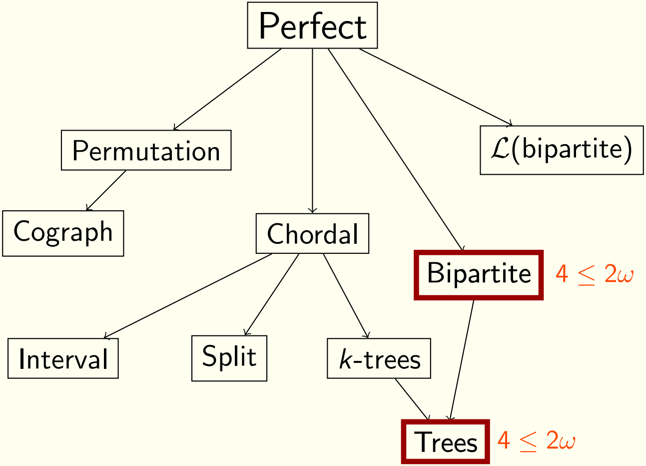
Link with 2-coloring of hypergraph



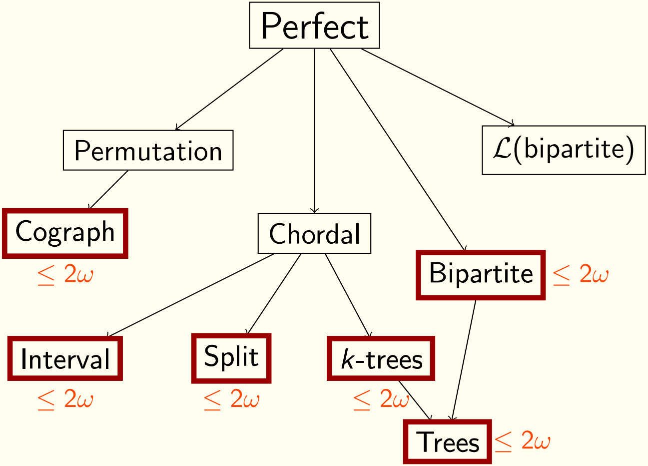
3-lid-coloring in bipartite graph $\Leftrightarrow$ 2-coloring in hypergraph

- $\rightarrow$ 3-Lid-Coloring in bipartite graph is NP-complete.
- $\rightarrow$ Polynomial in bipartite planar graphs with maximum degree 3.
- $\rightarrow$ k-regular bipartite graphs are 3 lid-colorable if $k \geq 4$.

Perfect graphs



Perfect graphs



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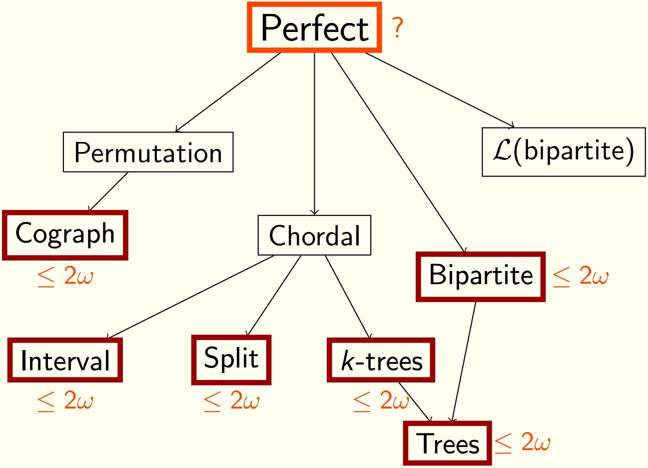
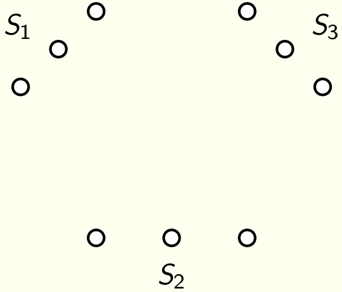
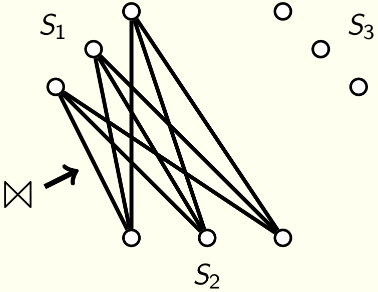


Perfect graphs are not anymore perfect



- S_1, S_2, S_3 stable set of size k.

Perfect graphs are not anymore perfect



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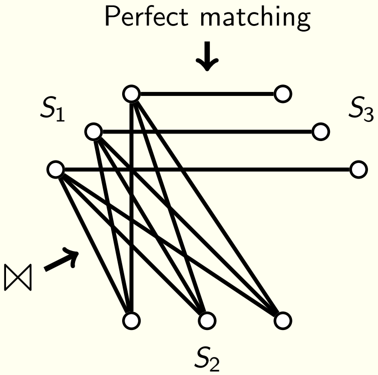

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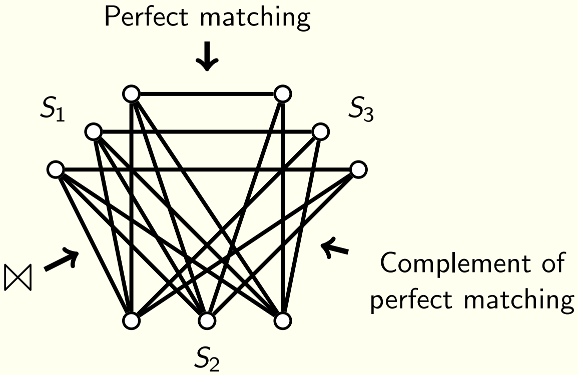

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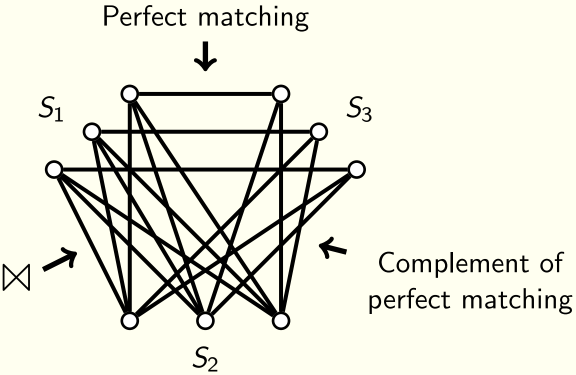
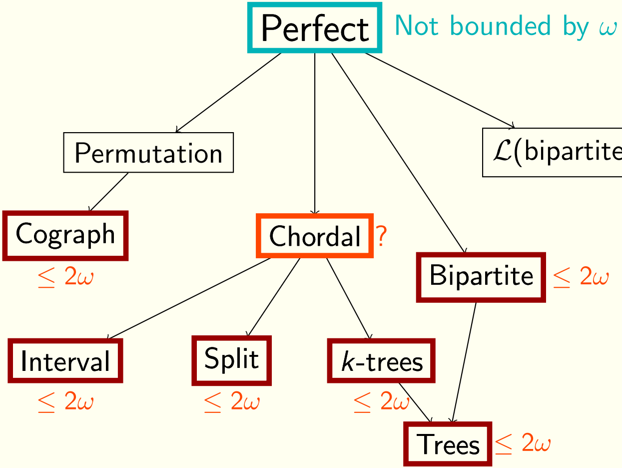


- S_1, S_2, S_3 stable set of size k.
- All vertices in S_2 must have different colors.
- $\chi_{lid} \geq k$, $\omega = 3$.

Perfect graphs - a conjecture



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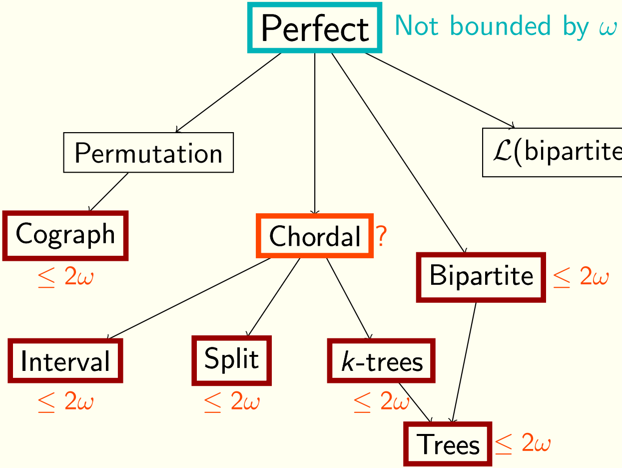
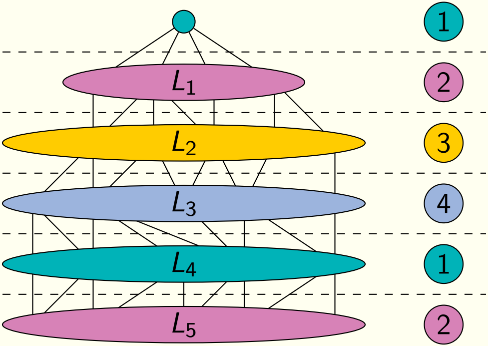


Conjecture Esperet, Gravier, Montassier, Ochem, P., 2012

Any chordal graph G has a lid-coloring with $2\omega(G)$ colors.

A good method for coloring



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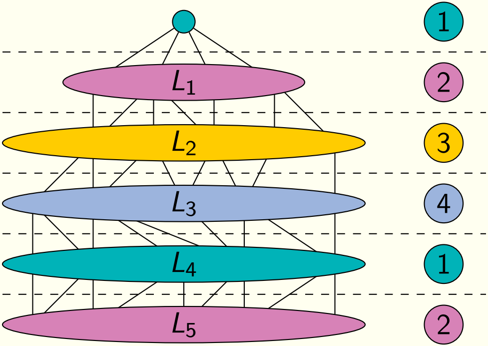

- Outerplanar graphs: L_i = union of paths, 5 colors
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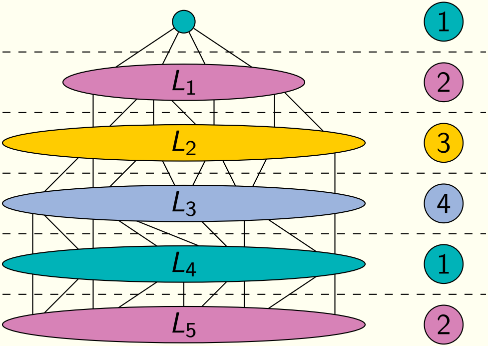

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- Planar graphs: L_i = outerplanar, 20 colors and 16 more colors
  $\rightarrow 4 \times 20 \times 16 = 1280$ colors (Gonçalves, P., Pinlou, 2012)

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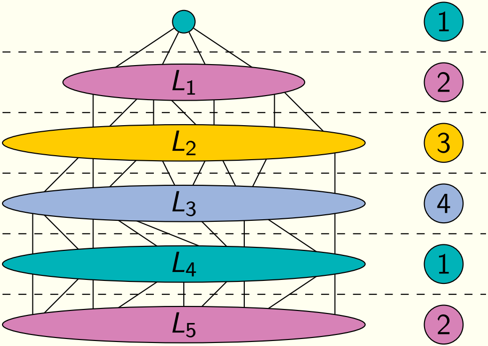

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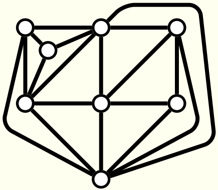
- Same idea for K_k-minor free graphs (Gonçalves, P., Pinlou, 2012)

Perspectives on identifying colorings

- Solve conjecture on chordal graphs.
- Better bound on planar graphs.



Worst example
8 colors

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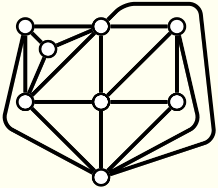

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- Tight bound with maximum degree: $\chi_{lid} \leq \mathbf{?}\Delta^2$

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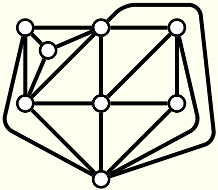


Worst example
8 colors

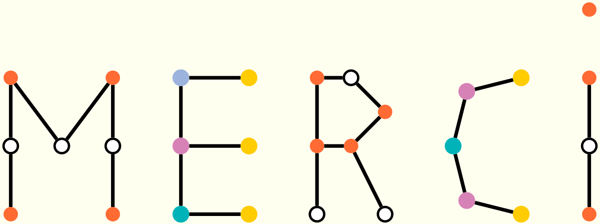
- Tight bound with maximum degree: $\chi_{lid} \leq$ **?** Δ^2
- Global version

Final conclusion and perspectives

- Bounds and extremal graphs
  $\rightarrow$ Conjecture Foucaud, Klasing, Kosowski, Raspaud

- Study in restricted classes of graphs
  $\rightarrow$ Other classes with $\gamma^{ID}(G) \geq \Theta(\sqrt{|V|})$?
  $\rightarrow$ Study in king grid, Sierpiński graphs, interval graphs
  $\rightarrow$ Is IDENTIFYING CODE polynomial for interval graphs ?

- Variations
  $\rightarrow$ Open questions on identifying colorings
  $\rightarrow$ Two other variations: weak and light codes, tolerant identifying codes
  $\rightarrow$ Generalization to hypergraph ?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